# CS 61C Fall 2024

## **RISC-V** Calling Convention

Discussion 4

### 1 Pre-Check

This section is designed as a conceptual check for you to determine if you conceptually understand and have any misconceptions about this topic. Please answer true/false to the following questions, and include an explanation:

- 1.1 After calling a function and having that function return, the t registers may have been changed during the execution of the function, while a registers cannot.
- 1.2 In order to use the saved registers (s0-s11) in a function, we must store their values before using them and restore their values before returning.
- 1.3 The stack should only be manipulated at the beginning and end of functions, where the callee saved registers are temporarily saved.

#### 2 Calling Convention

Let's review what special meaning we assign to each type of register in RISC-V.

Register	Convention	Saver
x0	Stores <b>zero</b>	N/A
sp	Stores the <b>stack pointer</b>	Callee
ra	Stores the <b>return address</b>	Caller
a0 - a7	Stores <b>arguments</b> and <b>return</b>	Caller
	values	
t0 - t6	Stores <b>temporary</b> values that $do$	Caller
	not persist after function calls	
s0 - s11	Stores <b>saved</b> values that $persist$	Callee
	after function calls	

To save and recall values in registers, we use the sw and lw instructions to save and load words to and from memory, and we typically organize our functions as follows:

```
# Prologue
1
   addi sp sp -8 # Room for two registers. (Why?)
2
    sw s0 0(sp) # Save s0 (or any saved register)
3
    sw s1 4(sp) # Save s1 (or any saved register)
4
5
    # Code omitted
6
7
    # Epilogue
8
9
   lw s0 0(sp) #Load s0 (or any saved register)
10
   lw s1 4(sp) #Load s1 (or any saved register)
11
    addi sp sp 8 #Restore the stack pointer
12
```

Now, let's see what happens if we ignore calling convention.

2.1 Consider the following blocks of code:

1	main:	1	foo:
2	# Prologue	2	# Preamble
3	# Saves ra	3	# Saves s0
4		4	
5	# Code omitted	5	# Code omitted
6	addi s0 x0 5	6	addi s0 x0 4
7	<b>#</b> Breakpoint 1	7	<b>#</b> Breakpoint 2
8	jal ra foo	8	
9	<b>#</b> Breakpoint 3	9	# Epilogue
10	mul a0 a0 s0	10	<b>#</b> Restores s0
11	# Code omitted	11	jr ra
12			
13	# Epilogue		
14	<b>#</b> Restores ra		
	i ovit		

15 jexit

(a) Does main always behave as expected, as long as foo follows calling convention?

- (b) What does s0 store at breakpoint 1? Breakpoint 2? Breakpoint 3?
- (c) Now suppose that foo didn't have a prologue or epilogue. What would s0 store at each of the breakpoints? Would this cause errors in our code?

In part (c) above, we saw one way how not following calling convention could make our code misbehave. Other things to watch out for are: assuming that a or t registers will be the same after calling a function, and forgetting to save ra before calling a function.

Function myfunc takes in two arguments: a0, a1. The return value is stored in a0. In myfunc, generate\_random is called. It takes in 0 arguments and stores its return value in a0.

```
myfunc:
1
         # Prologue (omitted)
2
3
         addi t0 x0 1
4
         slli t1 t0 2
5
         add t1 a0 t1
6
         add s0 a1 x0
7
8
         jal generate_random
9
10
         add t1 t1 a0
11
         add a0 t1 s0
12
13
         # Epilogue (omitted)
14
         ret
15
```

2.2 Which registers, if any, need to be saved on the stack in the prologue?

2.3 Which registers do we need to save on the stack before calling generate\_random?

2.4 Which registers need to be recovered in the epilogue before returning?

#### 3 Recursive Calling Convention

Write a function  $sum_square$  in RISC-V that, when given an integer n, returns the summation below. If n is not positive, then the function returns 0.

$$n^{2} + (n-1)^{2} + (n-2)^{2} + \ldots + 1^{2}$$

To implement this, we will use a tail-recursive algorithm that uses the a1 register to help with recursion. More specifically, you will be writing the following function:

sum_squares_recursive: Return the value $m + n^2 + (n-1)^2 + (n-2)^2 + \ldots + 1^2$					
Anorementa	a0	A 32-bit number. n. You may assume $n \leq 10000$ .			
Arguments	a1	A 32-bit number. m.			
Return value	a0	$m + n^2 + (n-1)^2 + (n-2)^2 + \ldots + 1^2$ . If $n \le 0$ , return $m$ .			

When the above function is called with a1 set to 0, we will get the behavior that we expect. For this problem, you are given a RISC-V function called square that takes in a single integer and returns its square.

square: Square a number					
Arguments	a0	n.			
Return value	a0	$n^2$			

3.1 Since this is a recursive function, let's start with the base case of our recursion.

sum\_squares:

\_\_\_\_\_ zero\_case

# To be implemented in the next question.

zero\_case:

----jr ra

3.2 Next, implement the recursive logic. Hint: If you let  $m' = m + n^2$ , then

$$m + n^{2} + (n - 1)^{2} + \ldots + 1^{2} = m' + (n - 1)^{2} + \ldots + 1^{2}$$

sum\_squares:

```
# Handle zero case (previous question)
mv t0 a0
jal ra _____
add a1 a0 a1
addi a0 t0 -1
jal ra ______
jr ra
zero_case:
    # Handle zero case (previous question)
jr ra
```

- 3.3 Now, think about calling convention from the caller perspective. After the call to square, what is in a0? a1? Which one of the registers will cause a calling convention violation?
- 3.4 What about the recursive call? What will be in a0 after the call to sum\_squares? a1?
- 3.5 Now, go back and fix the calling convention issues you identified. Note that not all blank lines may be used. There may also be another caller saved register that you need to save as well!

```
sum_squares:
```

jr ra

```
# Handle zero case (previous question)
  mv t0 a0
     _____
   _____
  # (previous question)
  jal ra _____
    _____
  _____
     _____
  add a1 a0 a1
  addi a0 t0 -1
    _____
  # (previous question)
  jal ra _____
   _____
    _____
  _____
zero_case:
  # Handle zero case (previous question)
```

#### 6 RISC-V Calling Convention

**3.6** Now, from a callee perspective, do we have to save any registers in the prologue and epilogue? If yes, what registers do we have to save, and where do we place the prologue and epilogue? If no, briefly explain why.