CS 61C Fall 2024

Instruction Translation, CALL

Discussion 5

1 Pre-Check

This section is designed as a conceptual check for you to determine if you conceptually understand and have any misconceptions about this topic. Please answer true/false to the following questions, and include an explanation:

- 1.1 The compiler may output pseudoinstructions.
- 1.2 The main job of the assembler is to perform optimizations on the assembly code.
- 1.3 The object files produced by the assembler are only moved, not edited, by the linker.
- 1.4 The destination of all jump instructions is completely determined after linking.

2 Translation

 $\fbox{2.1}$ In this question, we will be translating between RISC-V code and binary/hexadecimal values.

Translate the following RISC-V instructions into binary and hexadecimal notations.

- 1 addi s1 x0 -24 = 0b_____ = 0x_____
- 2 sh s1 4(t1) = 0b_____ = 0x_____

[2.2] In this question, we will be translating between RISC-V code and binary/hexadecimal values.

Translate the following hexadecimal values into the relevant RISC-V instruction.

- 1 0x234554B7 = _____
- 2 0xFE050CE3 = ______

3 RISC-V Addressing

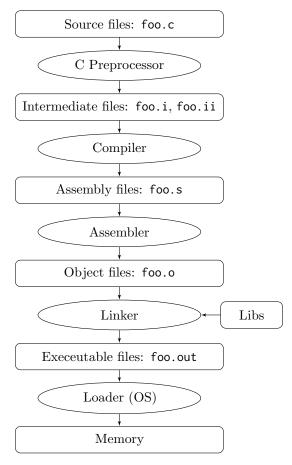
We have several addressing modes to access memory (immediate not listed):

- 1. Base displacement addressing adds an immediate to a register value to create a data memory address (used for lw, lb, sw, sb).
- 2. PC-relative addressing uses the PC and adds the immediate value of the instruction (multiplied by 2) to create an instruction address (used by branch and jump instructions).
- 3. Register Addressing uses the value in a register as an instruction address. For instance, jalr, jr, and ret, where jr and ret are just pseudoinstructions that get converted to jalr.
- 3.1 What is the range of 32-bit instructions that can be reached from the current PC using a branch instruction? Recall that RISC-V supports 16b instructions via an extension.
- 3.2 What is the maximum range of 32-bit instructions that can be reached from the current PC using a jump instruction?
- 3.3 Given the following RISC-V code (and instruction addresses), fill in the blank fields for the following instructions (you'll need your RISC-V reference sheet!). Each field refers to a different block of the instruction encoding.

1	0x002cff00: loop:	add t1, t2, t0	0x33
2	0x002cff04:	jal ra, foo	0x6F
3	0x002cff08:	bne t1, zero, loop	0x63
4			
5	0x002cff2c: foo:	jr ra	ra =

4 CALL

The following is a diagram of the CALL stack detailing how C programs are built and executed by machines.



- [4.1] How many passes through the code does the Assembler have to make? Why?
- [4.2] Which step in CALL resolves relative addressing? Absolute addressing?
- 4.3 Describe the six main parts of the object files outputted by the Assembler (Header, Text, Data, Relocation Table, Symbol Table, Debugging Information).