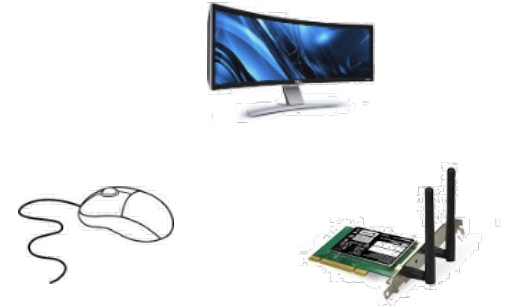
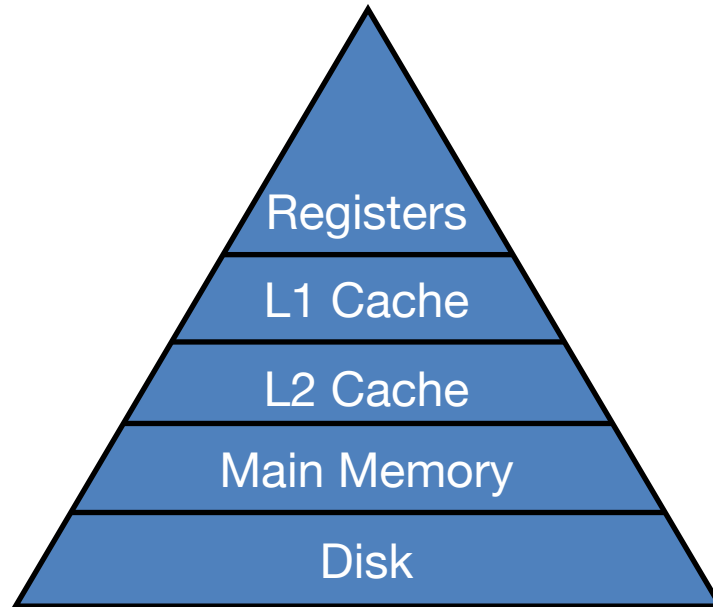


Operating Systems

Mid-Semester Survey

Operating Systems

Processor

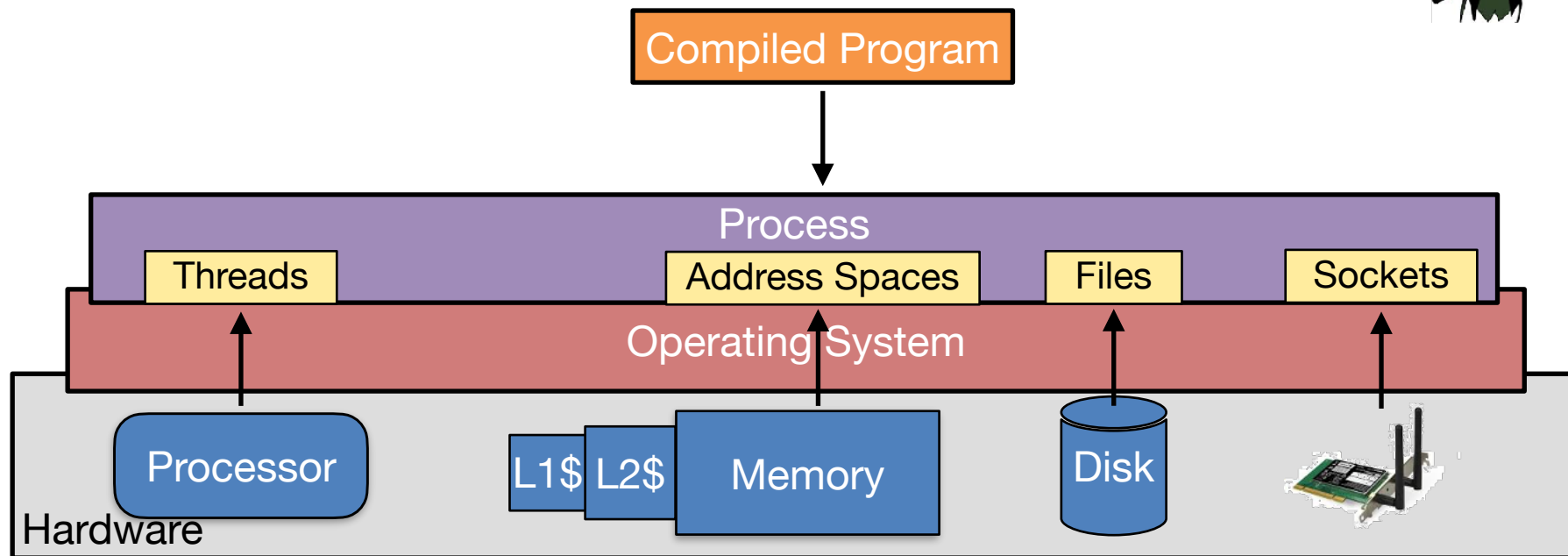


Operating Systems

- Unix
 - Berkeley Software Distribution (BSD)
 - macOS
- Linux distribution
 - Debian
 - Ubuntu
 - Red Hat
- Microsoft Windows

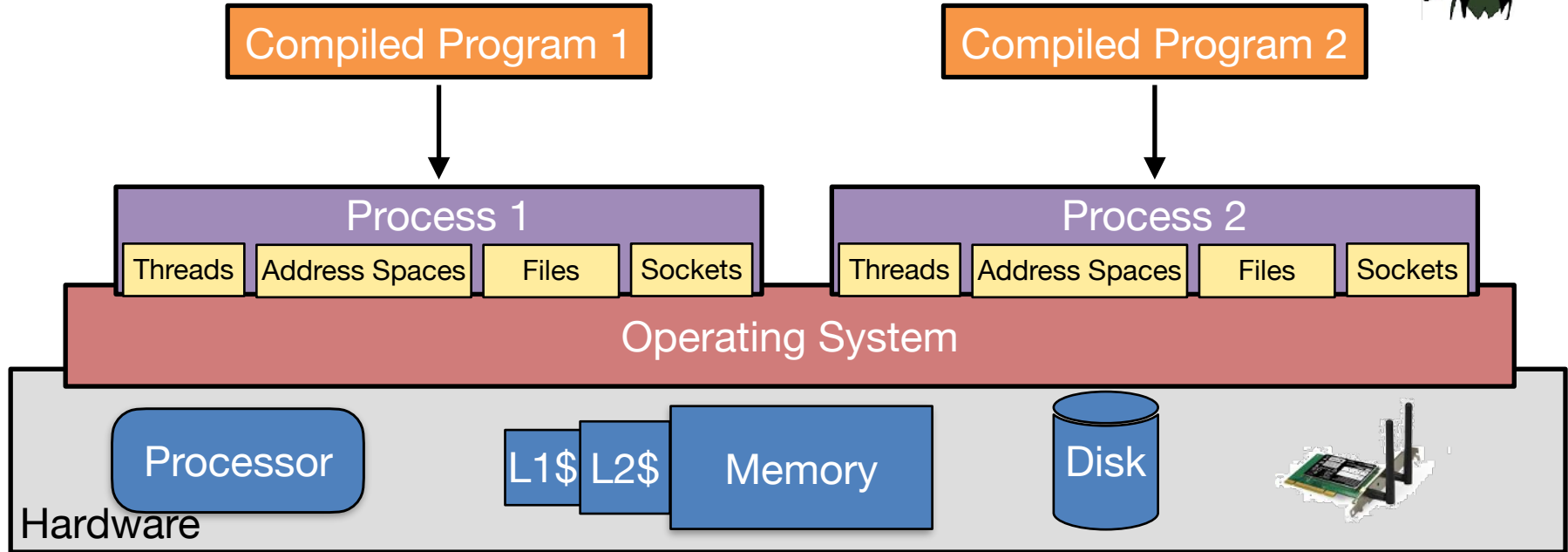
What is an Operating System

- Illusionist
 - Provide clean, easy-to-use abstractions of physical resources



What is an Operating System

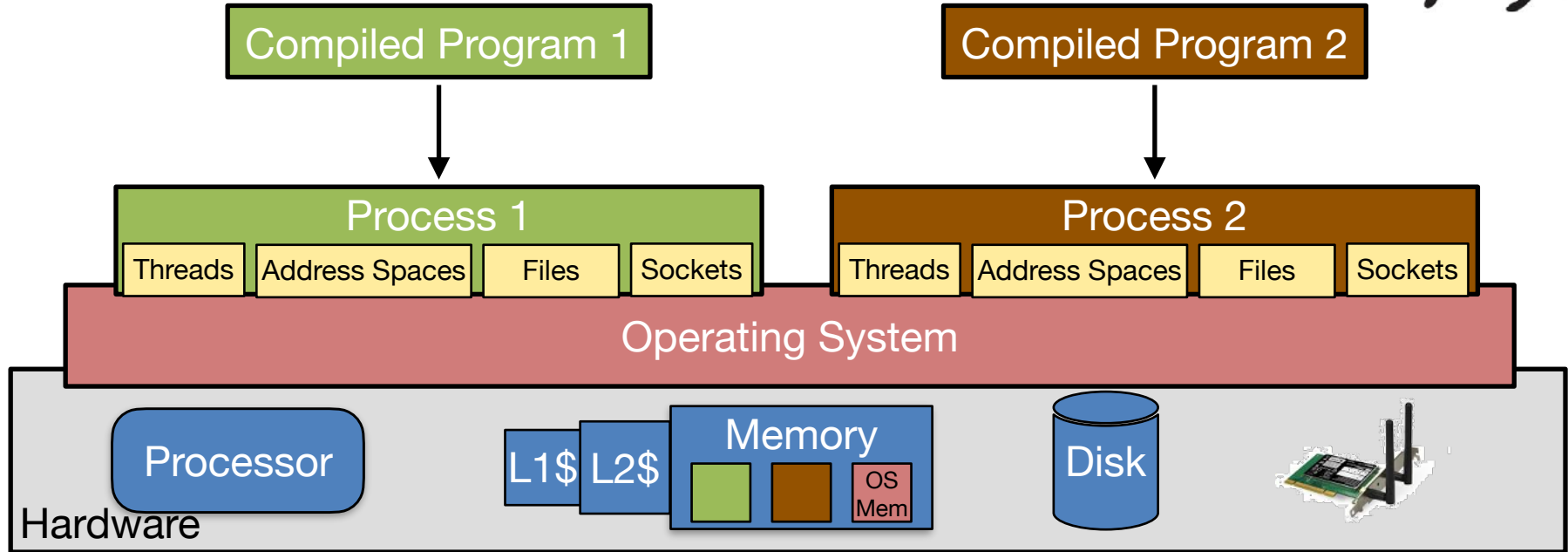
- Illusionist
 - Provide clean, easy-to-use abstractions of physical resources



What is an Operating System

- Referee

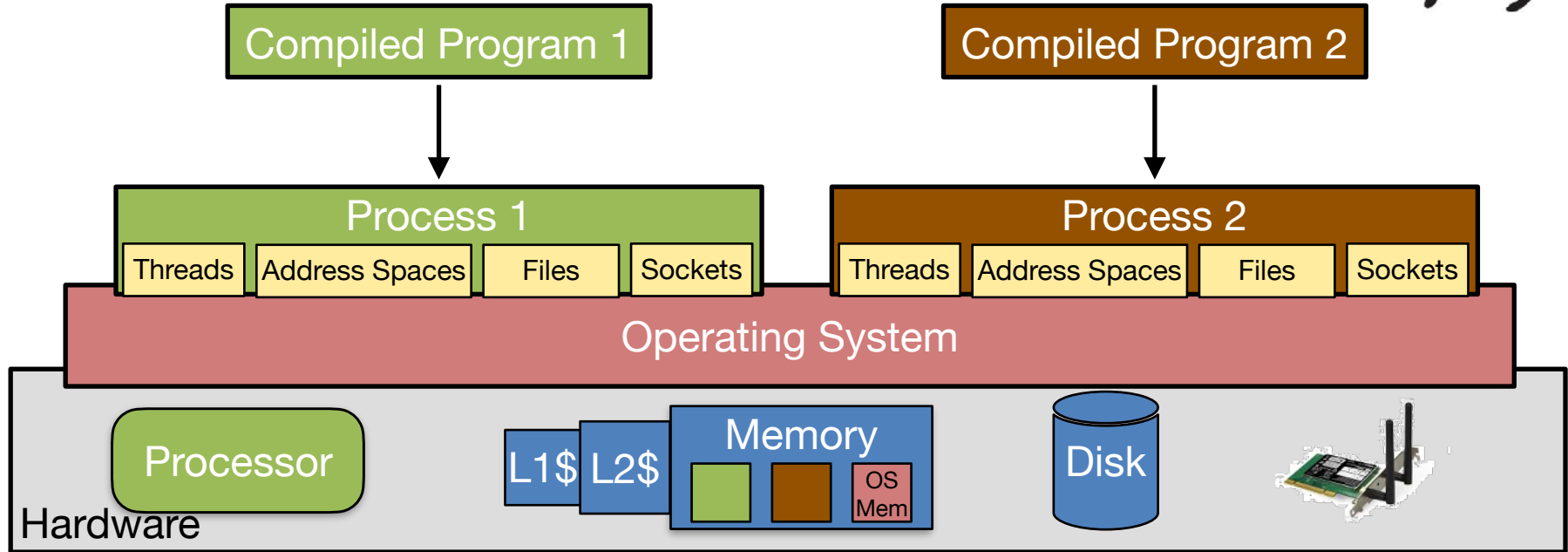
- Manage protection, isolation, and sharing of resources



What is an Operating System

- Referee

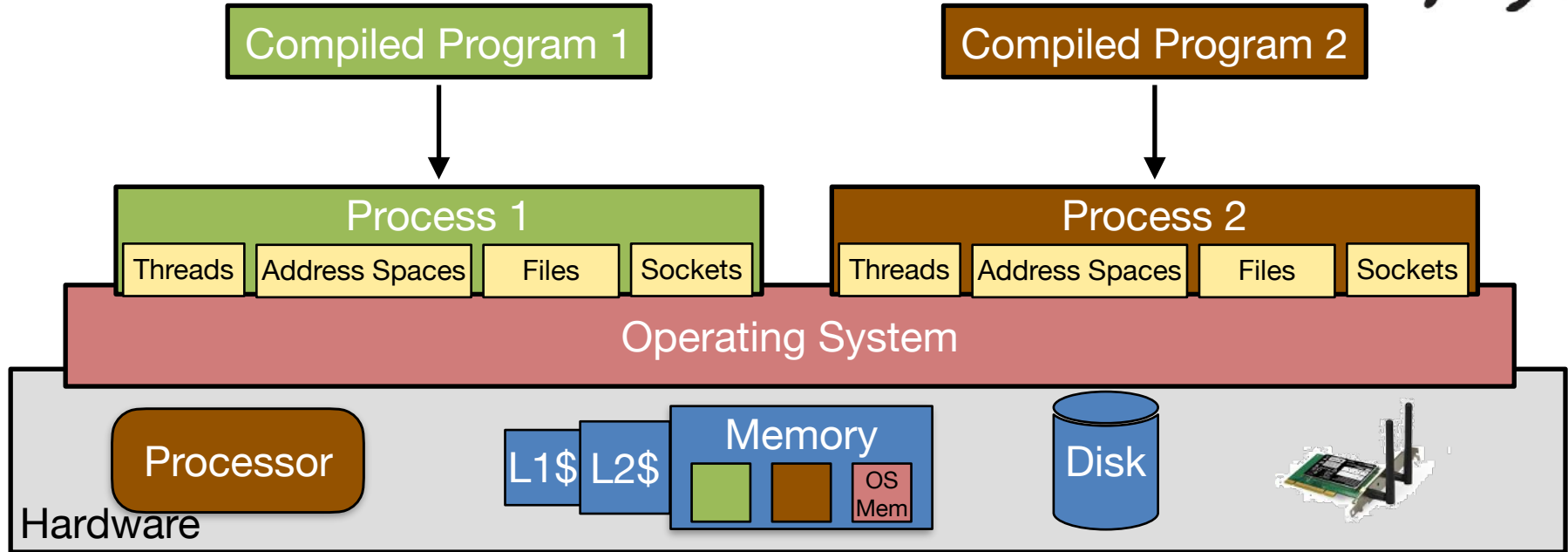
- Manage protection, isolation, and sharing of resources



What is an Operating System

- Referee

- Manage protection, isolation, and sharing of resources



Context Switch

- Context switch: Switching from executing one program to another
- Allows multiple processes to run on the same processor
- The OS determines when to context switch

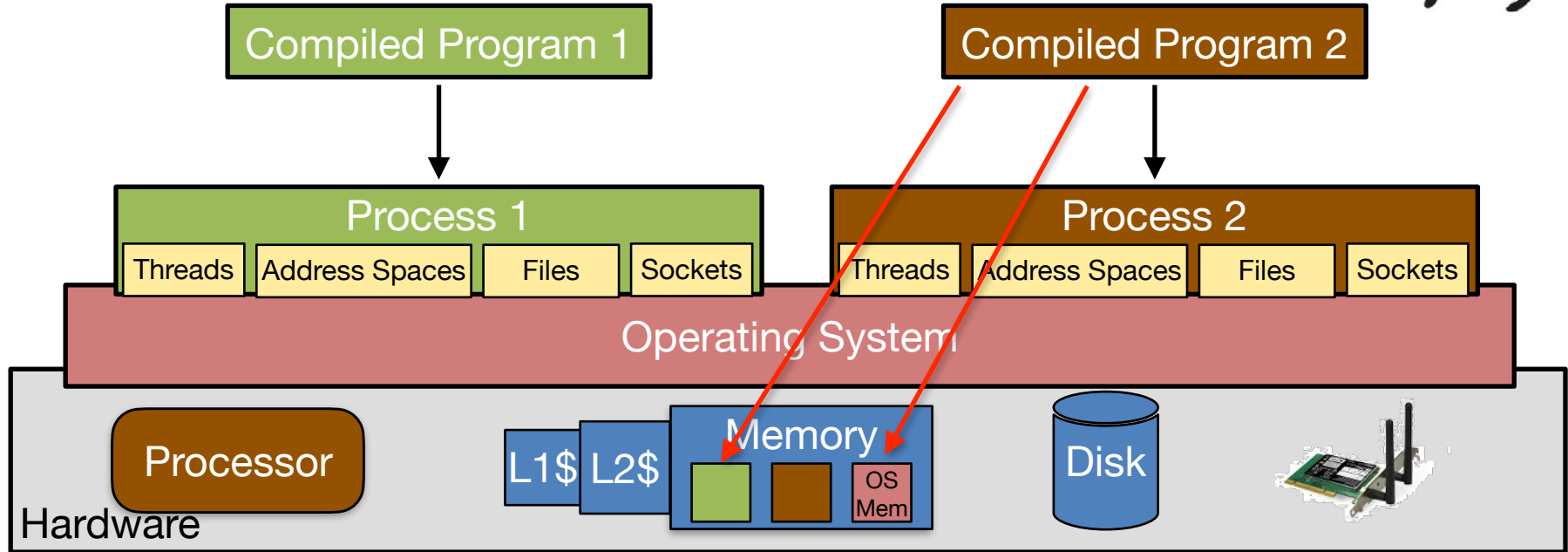
What happens on a context switch?

1. The OS takes control of the CPU from the current process
2. The OS saves the state of the current process
3. The OS loads the state of the next process
4. The OS hands over the CPU to the next process

What is an Operating System

- Referee

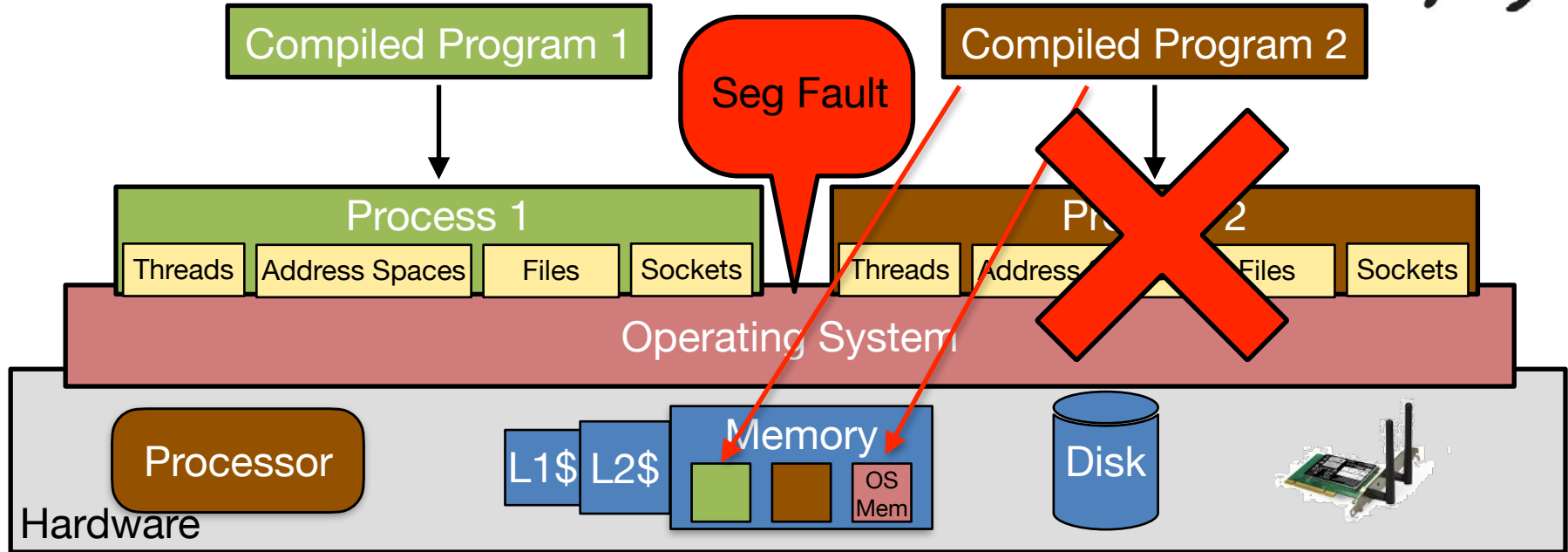
- Manage protection, isolation, and sharing of resources



What is an Operating System

- Referee

- Manage protection, isolation, and sharing of resources



Protection

- OS isolates processes from each other
- OS isolates itself from other processes
- ... even though they are actually running on the same hardware!

Dual Mode Operation

- Hardware provides at least two modes:
 1. Kernel Mode (or “supervisor” mode)
 2. User Mode
- Certain operations are prohibited when running in user mode
 - interacting directly w/ hardware, writing to kernel memory
- OS mostly runs in user mode
- Switching between user mode and kernel mode
 - System calls, interrupts, exceptions

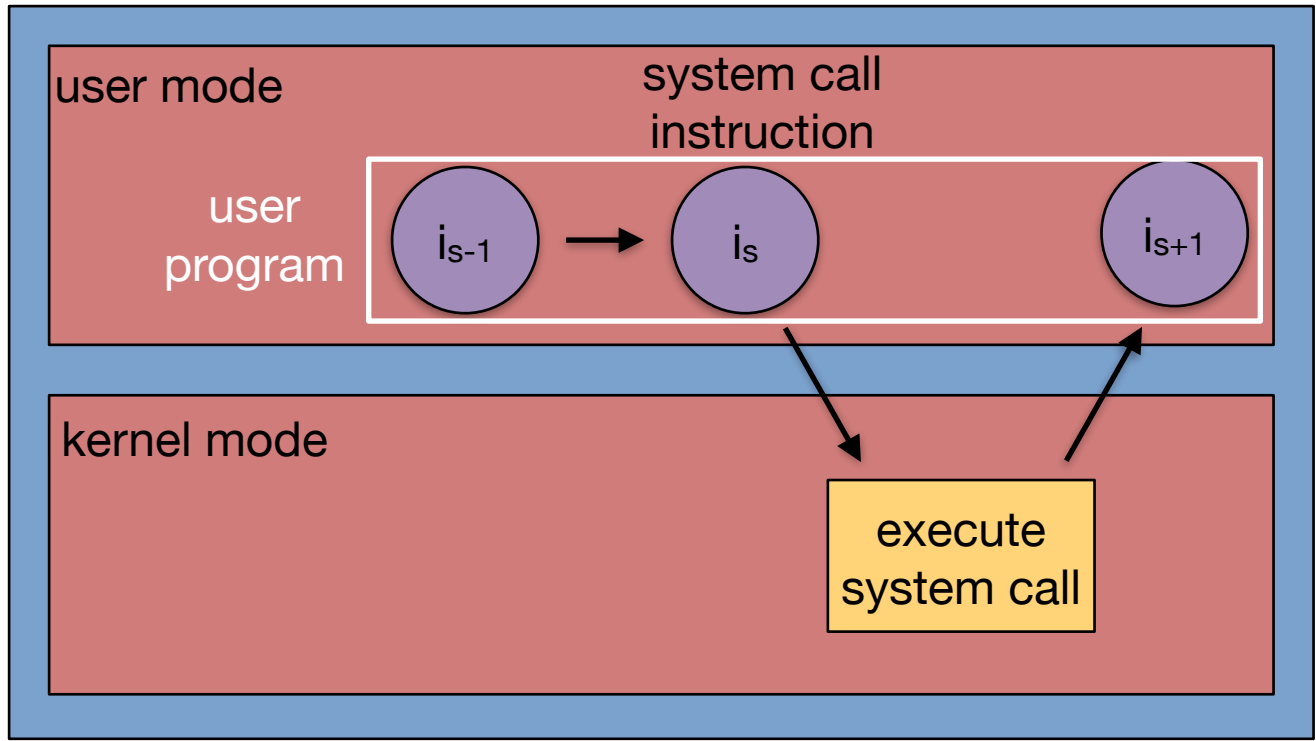
System Calls (syscall)

- Allows the program to request a service from the operating system
- Examples
 - creating and deleting files
 - reading and writing files
 - accessing external devices like a scanner
 - (ecalls in RISC-V)
- Similar to function calls except it's executed by the kernel

System Calls

● Single instruction

■ Series of Instructions



Interrupts vs Exceptions

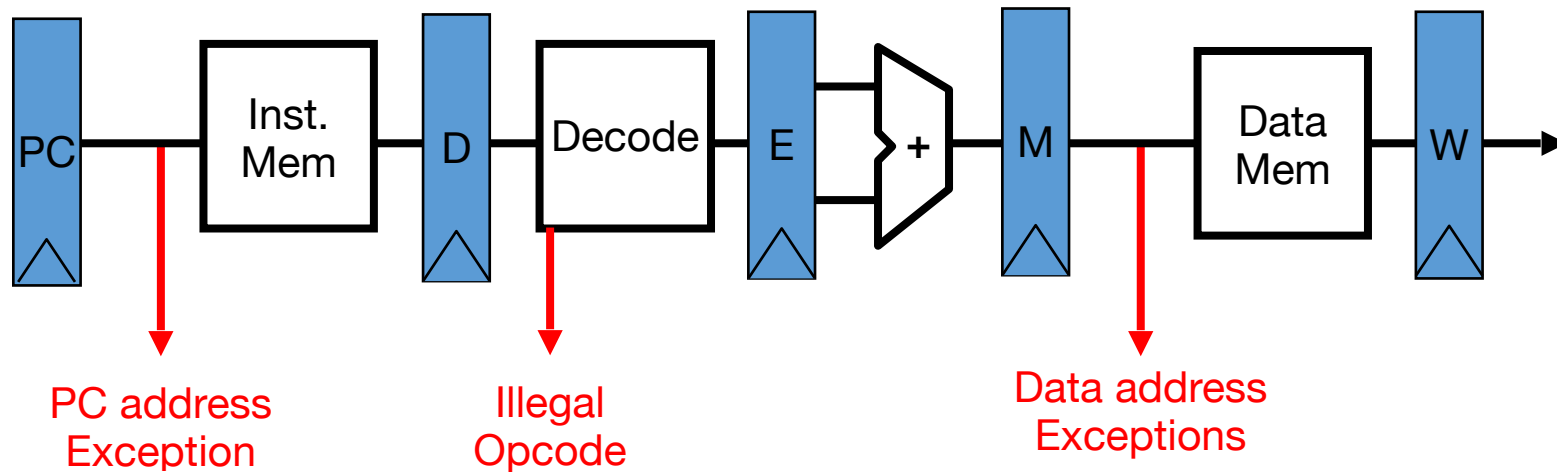
- Interrupts

- Caused by an event *external* to the current running program
- Ex: Key press
- Asynchronous to the current program
 - Does not need to be handled immediately, but should be handled soon

- Exceptions

- Caused by an event *during* the execution of the current program
- Ex: illegal instruction, divide by zero
- Synchronous
 - Must be handled immediately

Exceptions



How to Handle Interrupts and Exceptions?

- Trap Handler: code that services the interrupt or exception
- From the program's point of view, it must look like nothing happened

Traps

1. All instructions before the faulting instructions must complete
2. All instructions after to the faulting instruction must be flushed
3. The faulting instruction must be flushed
4. Execution of the trap handler begins

What does the Trap Handler Do?

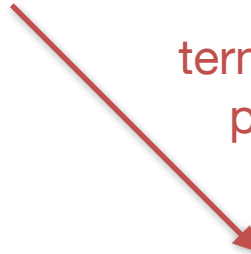
1. Save the state of the current program
 - Save ALL of the registers
2. Determine what caused the exception or interrupt
3. Handle exception or interrupt

continue execution
of the program



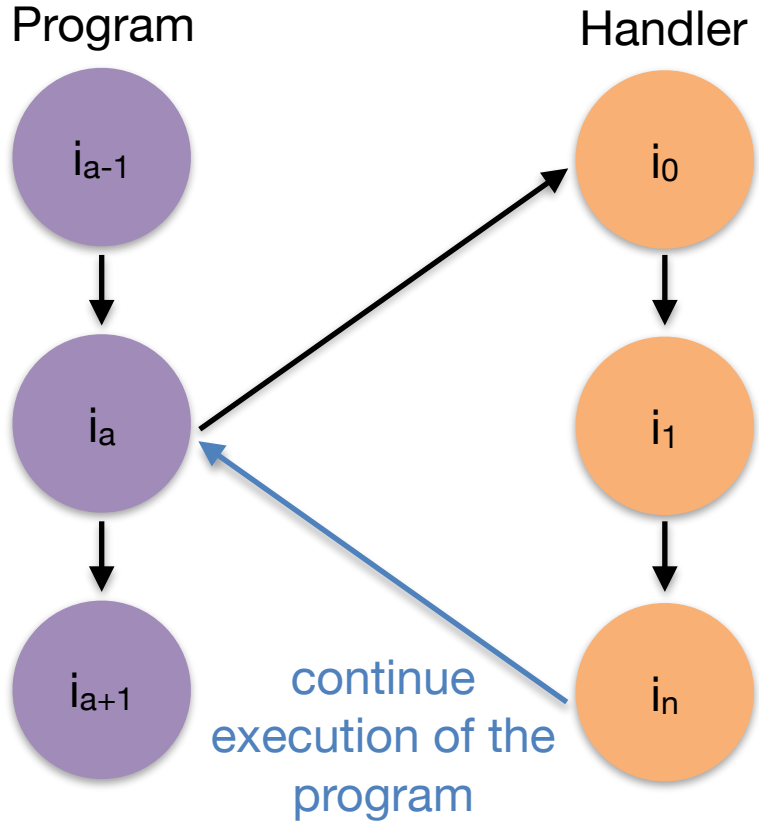
4. Restore the state of the program
5. Return control to the program

terminate the
program



4. Terminate the program (free resources, etc)
5. Schedule a new program

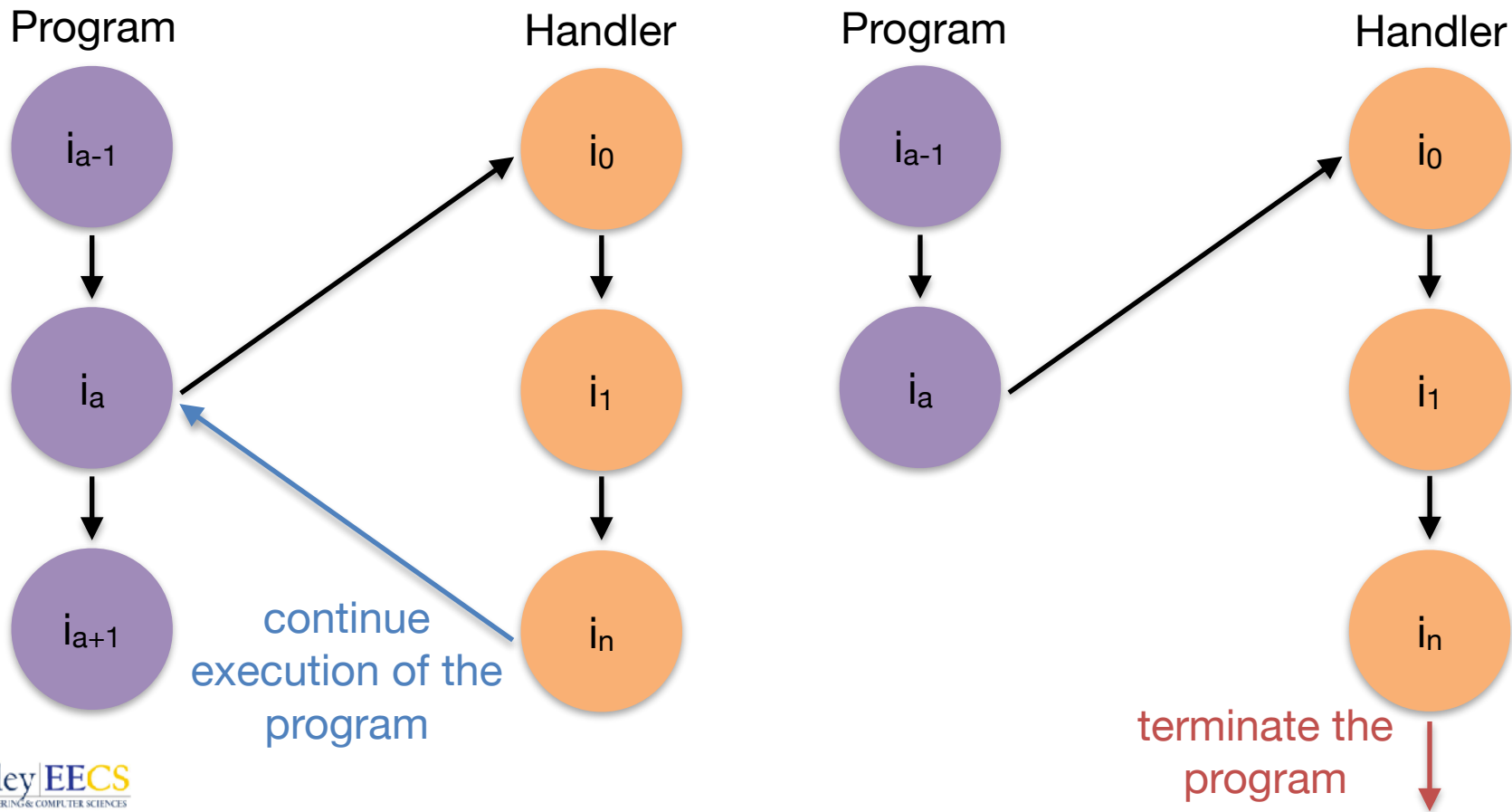
Traps



i_a is the instruction that caused the exception

The program continues at the instruction that caused the exception

Traps



Which path to choose?

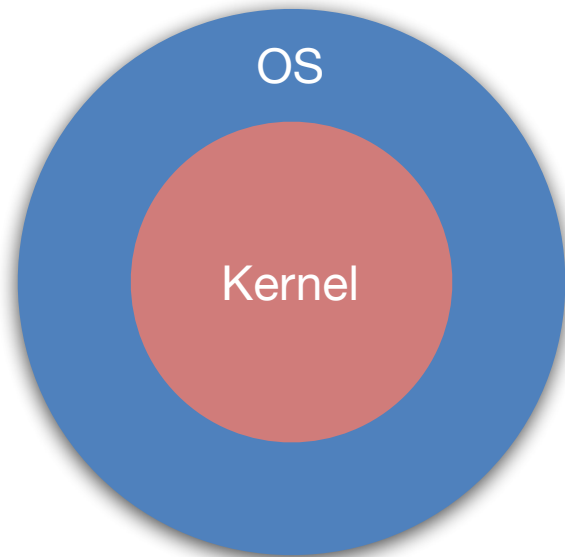
- Continue execution of the program
 - Interrupts (most likely)
 - Certain memory exceptions (we'll see more later)
- Terminate program
 - Illegal instruction
 - Certain illegal memory accesses

Program's Point of View

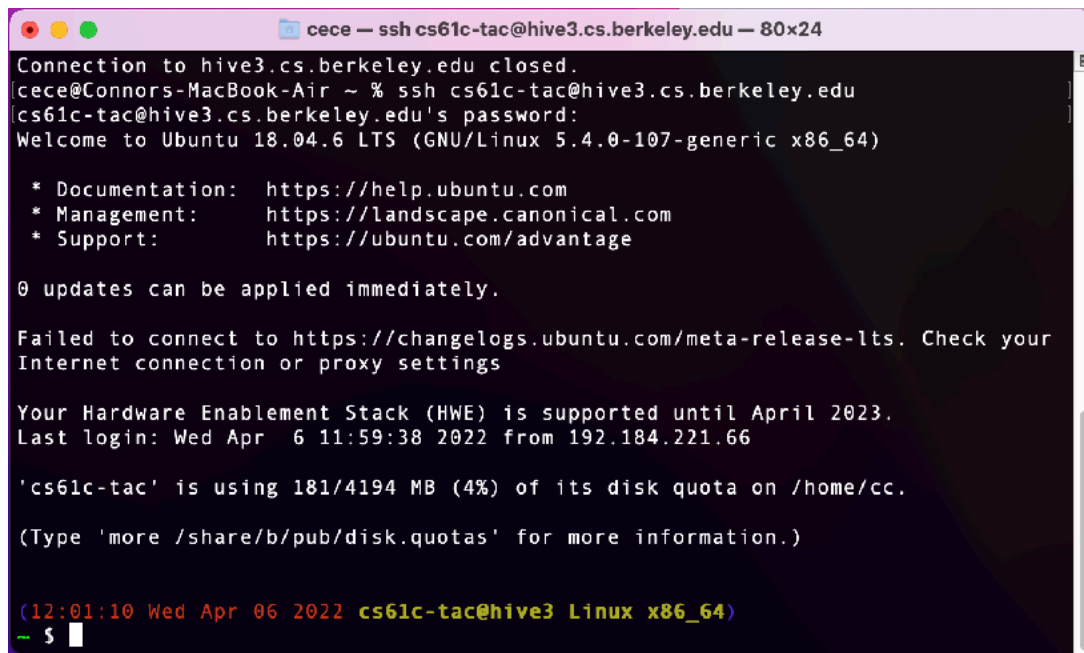
- Almost like nothing ever happened
 - The program state did not change
- Maybe large gap in between one instruction and the next
- Caches may have been trashed
 - Because something else was using them

Kernel

- Core of the OS
- Manages resources
 - scheduling, memory, I/O
- Things not in the kernel
 - User interface
 - Networking
- Lots of variation between different operating systems



- A program that exposes the operating system's services



```
cece — ssh cs61c-tac@hive3.cs.berkeley.edu — 80x24
Connection to hive3.cs.berkeley.edu closed.
cece@Connors-MacBook-Air ~ % ssh cs61c-tac@hive3.cs.berkeley.edu
cs61c-tac@hive3.cs.berkeley.edu's password:
Welcome to Ubuntu 18.04.6 LTS (GNU/Linux 5.4.0-107-generic x86_64)

 * Documentation:  https://help.ubuntu.com
 * Management:    https://landscape.canonical.com
 * Support:       https://ubuntu.com/advantage

0 updates can be applied immediately.

Failed to connect to https://changelogs.ubuntu.com/meta-release-lts. Check your
Internet connection or proxy settings

Your Hardware Enablement Stack (HWE) is supported until April 2023.
Last login: Wed Apr  6 11:59:38 2022 from 192.184.221.66

'cs61c-tac' is using 181/4194 MB (4%) of its disk quota on /home/cc.
(Type 'more /share/b/pub/disk.quotas' for more information.)

(12:01:10 Wed Apr 06 2022 cs61c-tac@hive3 Linux x86_64)
-- $
```

What happens at Boot?

1. The BIOS (Basic Input/Output System) runs
 - Power-on-self-test (POST)
 - The BIOS finds and executes the bootloader
2. The bootloader loads in part of the operating system
3. The operating system initializes services, drivers, etc
4. Launch a process that waits for an input in a loop

Bootstrapping: A chain of stages, in which at each stage, a smaller, simpler program loads and then executes the larger, more complicated program of the next stage (Wikipedia)

How do you begin executing a program?

- Loader: responsible for loading programs into memory
 1. The loader loads program into memory
 2. The loader sets `argc` and `argv`
 3. The OS jumps to `main` and transfers control to the process

What is an Operating System?

- Provides *isolation* between running processes
 - Each program runs in its own world
- Provides *interaction* with the outside world
 - interact with devices like mouse, display, network

What is an Operating System

- Illusionist
 - Provide clean, easy-to-use abstractions of physical resources
 - Masks limitations
 - Higher level objects: files, sockets
- Referee
 - Manage protection, isolation, and sharing of resources
 - Resource allocation and communication



Coming up...

- Virtual memory
 - How the OS isolates processes
- I/O
 - How the OS communicates with the outside world